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10/534,873	11/22/2005	Marc Joye	032326-301	6724
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EXAMINER				
WRIGHT, BRYAN F				
ART UNIT		PAPER NUMBER		
2431				
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**Please find below and/or attached an Office communication concerning this application or proceeding.**

The time period for reply, if any, is set in the attached communication.

Notice of the Office communication was sent electronically on above-indicated "Notification Date" to the following e-mail address(es):

ADIPFDD@bipc.com

### Office Action Summary

**Application No.**

10/534,873

**Applicant(s)**

JOYE ET AL.

**Examiner**

BRYAN WRIGHT

**Art Unit**

2431

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --  
**Period for Reply**

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

**Status**

- 1) ☒ Responsive to communication(s) filed on 22 November 2005.
- 2a) ☐ This action is **FINAL**. 2b) ☒ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

**Disposition of Claims**

- 4) ☒ Claim(s) 1-13 is/are pending in the application.
- 4a) Of the above claim(s) \_\_\_\_\_ is/are withdrawn from consideration.
- 5) ☐ Claim(s) \_\_\_\_\_ is/are allowed.
- 6) ☒ Claim(s) 1-13 is/are rejected.
- 7) ☐ Claim(s) \_\_\_\_\_ is/are objected to.
- 8) ☐ Claim(s) \_\_\_\_\_ are subject to restriction and/or election requirement.

**Application Papers**

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on \_\_\_\_\_ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.  
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).  
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

**Priority under 35 U.S.C. § 119**

- 12) ☒ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☒ All b) ☐ Some \* c) ☐ None of:
1. ☒ Certified copies of the priority documents have been received.
  2. ☐ Certified copies of the priority documents have been received in Application No. \_\_\_\_\_.
  3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

\* See the attached detailed Office action for a list of the certified copies not received.

**Attachment(s)**

- 1) ☒ Notice of References Cited (PTO-892)
- 2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)
- 3) ☒ Information Disclosure Statement(s) (PTO-893)  
Paper No(s)/Mail Date 5/12/2005

- 4) ☐ Interview Summary (PTO-413)  
Paper No(s)/Mail Date: \_\_\_\_\_
- 5) ☐ Notice of Informal Patent Application
- 6) ☐ Other: \_\_\_\_\_

**DETAILED ACTION**

1. This action is in response to the original filing on 11/22/2005. Claims 1-13 are pending.

***Priority***

2. Applicant's claim for benefit of foreign priority under 35 U.S.C. 119 (a) - (d) is acknowledged.

***Drawings***

The subject matter of this application admits of illustration by a drawing to facilitate understanding of the invention. Applicant is required to furnish a drawing under 37 CFR 1.81(c). No new matter may be introduced in the required drawing. Each drawing sheet submitted after the filing date of an application must be labeled in the top margin as either "Replacement Sheet" or "New Sheet" pursuant to 37 CFR 1.121(d).

***Claim Rejections - 35 USC § 101***

35 U.S.C. 101 reads as follows:

Whoever invents or discovers any new and useful process, machine, manufacture, or composition of matter, or any new and useful improvement thereof, may obtain a patent therefor, subject to the conditions and requirements of this title.

3. Claims 1-10 are rejected under 35 U.S.C. 101 because the claimed invention is directed to non-statutory subject matter. Applicant's claims 1-10 are directed towards

an algorithm method that can be performed via physical computation using a piece of paper and pencil. The Office's current position is that such claims involving a algorithm method with functional descriptive material, do not fall within any of the categories of patentable subject matter set forth in 35 U.S.C. § 101, and such claims are therefore ineligible for patent protection. Applicant is advised to amend claims to read, "A method of integer division performed on a processor within an electronic device comprising".

***Claim Rejections - 35 USC § 112***

The following is a quotation of the second paragraph of 35 U.S.C. 112:

The specification shall conclude with one or more claims particularly pointing out and distinctly claiming the subject matter which the applicant regards as his invention.

4. Claims 3, 6, 10 and 11 are rejected under 35 U.S.C. 112, second paragraph, as being indefinite for failing to particularly point out and distinctly claim the subject matter which applicant regards as the invention. The Examiner contends claims 3, 6, 10, and 11 are improper claims on the basis that the meets and bounds of the claims cannot be readily ascertained. For example, applicant's claims recites "Input :  $a = (0, a_{m-1} \dots, a_0)$ ". The Examiner contends that one of ordinary skill in the art cannot reasonably identify an appropriate input bounds. The Examiner maintains the same assertion for the remaining claim limitation elements. The applicant is advised to re-write the claim in a form that would allow a clear and concise understanding of the meets and bounds of the claim limitations. For example the claim may read, "an input consisting of integer value  $a$ ". Appropriate correction is required.

***Claim Rejections - 35 USC § 103***

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

This application currently names joint inventors. In considering patentability of the claims under 35 U.S.C. 103(a), the examiner presumes that the subject matter of the various claims was commonly owned at the time any inventions covered therein were made absent any evidence to the contrary. Applicant is advised of the obligation under 37 CFR 1.56 to point out the inventor and invention dates of each claim that was not commonly owned at the time a later invention was made in order for the examiner to consider the applicability of 35 U.S.C. 103(c) and potential 35 U.S.C. 102(e), (f) or (g) prior art under 35 U.S.C. 103(a).

5. Claims 1-13 are rejected under 35 U.S.C. 103(a) as being unpatentable over Menezes (NPL "Handbook of applied cryptography" (cited from IDS)) in view of Drexler et al. (US 2003/0061498 and Drexler hereinafter)

6. As to claim 1, Menezes teaches a cryptographic method during which an integer division of the type  $q = a \text{ div } b$  and  $r = a \text{ mod } b$  is performed (i.e., ...teaches integer division [pg. 63, sect. 2.79]), with where  $q$  is a quotient [pg. 64, sect. 2.82],  $a$  is a number of containing  $m$  bits [pg. 64, sect. 2.82],  $b$  is a number of containing  $n$  bits [pg.

64, sect. 2.82], with  $n$  less than or equal to  $m$  and  $b_{n-1}$  is non-zero,  $b_{nq}$  being the most significant bit of  $b$  [pg. 64, sect. 2.83], a comprising the following steps:

(i) performing a partial division of a word  $A$ , comprising left  $n$  bits of the number  $a$ , by the number  $b$  to obtain a bit of the quotient  $q$ , (i.e., ...teaches integer division [pg. 63, sect. 2.79]);

Menezes does not expressly teach:

(ii) repeating step (i) for  $m-n + 1$  iterations (e.g., For Loop) with the same operations being performed at each iteration, regardless of the value of the quotient bit obtained, to obtain the quotient  $q$ .

However, Menezes discloses instruction for which could be implemented as a "Computer For Loop Condition Statement" for iterative calculation of the encryption process as recited on [pg. 598, sect. 14.20]. Therefore given applicants minor change of the "For Loop" instruction to be iterated through (e.g., carryout by the computer) the encryption process, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Menezes's "For Loop Condition Statement" by employing the well known feature of adding an additional iterative step (e.g.,  $(n+1)$ ) for which will enhance data encryption within a chip card [pg. 598, sect. 14.20].

Menezes does not expressly teach:

wherein at least one of the numbers a and b comprises secret data and  
(iii) generating encrypted or decrypted data in accordance with said quotient.

However, these features are well known in the art and would have been an obvious modification of the system disclosed by Menezes as introduced by Drexler. Drexler discloses:

wherein at least one of the numbers a and b comprises secret data (to provide use of a secret key (e.g., secret data) in the calculation [par. 18]) and  
(iii) generating encrypted or decrypted data in accordance with said quotient (to provide a procedure for encrypting text with the use of a quotient [par. 23; fig. 1].

Therefore, given the teachings of Drexler, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Menezes by employing the well known feature of encrypting text using a quotient produced in integer division disclosed above by Drexler, for enhancing encryption in a chip card. [fig. 1].

7. As to claim 2, Menezes a method where at each iteration, an addition of the number b to the word A and a subtraction of the number b from the word A are performed [pg. 598, sect. 14.20, 3.1].

8. As to claim 3, Menezes a method where all the following steps are performed:

Input  $a = (0, a_{m-1}, \dots, a_0)$   $b = (b_{n-1}, \dots, b_0)$  [pg. 598, sect. 14.20],

Output:  $q = a \div b$  and  $r = a \bmod b$  [pg. 598, sect. 14.20].

Menezes does not expressly teach:

$A = (0, a_{m-1}, \dots, a_{m-n+0}; o' < -1$

For  $j = 1$  to  $(m-n+1)$ , do:

$a \leftarrow \text{SHL}_{m+1}(a, 1); o' \leftarrow \text{carry } A - (c') \text{SUB}, (A, b) + (\sim') \text{ADD}, (A, b) \text{ } o \sim \neg(o' \text{ AND } \sim') /$

$(o' \text{ AND } \text{carry}) / (o' \text{ AND } \text{carry})$

$\text{lsb}(a) \text{ } g' \sim' \neg(3'$

End For

if  $(\sim o = \text{TRUE})$  then  $A \leftarrow \text{ADD}_n(A, b)$ .

However, Menezes discloses instruction for performing encryption utilizing integer division which could be implemented in the form of a "Computer For Loop Condition Statement". Menezes' iterative calculation of the encryption process is recited on [pg. 598, sect. 14.20]. Therefore given applicants minor change of the "For Loop" instruction to be iterated through (e.g., carryout by the computer) the encryption process, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Menezes's "For Loop Condition Statement" by employing the well known feature of adding an additional iterative step (e.g.,  $(n+1)$ ) for which will enhance data encryption within a chip card [pg. 598, sect. 14.20].



9. As to claim 4, Menezes teaches a method where at each iteration (e.g., "For Loop Iteration"), either the number b or of a number complementary to the number b is added to the word A [pg. 598, sect. 14.20].

10. As to claim 5, Menezes teaches a method further at each iteration, an of updating of a first variable (c') (e.g., "x") indicating whether, during the following iteration, the number b or the number b must is to be added with the word A according to the quotient bit produced. [pg. 598, sect. 14.20].

11. As to claim 6, Menezes teaches a method where all the following steps are performed :

Input:  $a = (0, a_{m-1}, \dots, a_0)$   $b = (b_{-1}, \dots, b_0)$  [pg. 598, sect. 14.20], and

Output:  $q = a \text{ div } b$  and  $r = a \text{ mod } b$  [pg. 598, sect. 14.20].

Menezes does not expressly teach:

$A = (0, a_{m-1}, \dots, a_{m-1})$ ;  $o' \leftarrow 1$ ;  $b \leftarrow \text{CPL}2_n(b)$

For  $j = 1$  to  $(m-n+1)$ , do:  $a \leftarrow \text{SHL}_{m+1}(a, 1)$ ;  $o' \leftarrow \text{carry}$

$d_{\text{addr}} \leftarrow A - b_{\text{addr}} + o'(b' - b_{\text{addr}})(c')\text{SUB},(A, b) + (\sim')\text{ADD},(A, b) \text{ c} \sim -(\sim' \text{ AND } \sim') / (c'$

$\text{AND carry}) / (c' \text{ AND carry}) \text{ lsb}(a) \text{ g}' \sim' - (3'$

End For

if  $(\sim = \text{TRUE})$  then  $A \leftarrow \text{ADD},(A, b)$ .

However, Menezes discloses instruction for performing encryption utilizing integer division which could be implemented as a "Computer For Loop Condition Statement" Menezes' iterative calculation of the encryption process is recited on [pg. 598, sect. 14.20]. Therefore given applicants minor change of the "For Loop" instruction to be iterated through (e.g., carryout by the computer) the encryption process, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Meneze's "For Loop Condition Statement" by employing the well known feature of adding an additional iterative step (e.g.,  $(n+1)$ ) for which will enhance data encryption within a chip card [pg. 598, sect. 14.20].

12. As to claim 7, Menezes teaches a method during which further including the steps, at each iteration, of performing an operation of complement to  $2^n$  of an updated data item (b or b ) or of a notional data item (c or c) and then-an adding the updated data item with the word A (pg. 598, 14.20, lines 2 and 3).

13. As to claim 8, Menezes teaches a method, further including the step, at each iteration, of updating a second variable (e.g., "n") is-also indicating whether, during the following iteration, the operation of complement to  $2^n$  must is to be performed on the updated data item or on the notional data item. (pg. 598, 14.20, lines 2 and 3).

14. As to claim 9, Menezes teaches a method further including the step, at each iteration, of updating of a third variable (e.g., "x") indicating whether the updated data item is equal to the data item b or to its complement to  $2^n$  (pg. 598, sect. 14.20).

15. As to claim 10, Menezes teaches a method according to one-of which claim 7, wherein all the following steps are also performed:

Input  $a = (0, a_{m-1}, \dots, a_0)$   $b = (b_{n-1}, \dots, b_0)$  [pg. 598, sect. 14.20], and

Output:  $q = a \text{ div } b$  and  $r = a \text{ mod } b$  [pg. 598, sect. 14.20].

Menezes does not expressly teach:

$o' < -1$  ;  $B < -1$ ,  $y < -1$  ;  $A = (0, a_m, \dots, a_{m-n+1})$

for  $j = 1$  to  $(m-n+1)$ , do:

$a \leftarrow \text{SHL}_{m+1}(a, 1); a \leftarrow \text{carry} \ll 8 - o' \ll 13$

$d_{\text{addr}} \leftarrow b_{\text{addr}} + 8 \text{ (} C_{\text{addr}} \leftarrow b_{\text{addr}} \text{)}$

$d \leftarrow \text{CPL2}.(d)$

$A \leftarrow \text{ADDn}(A, b) \text{ } o \leftarrow (o \text{ AND } o') / (cr \text{ AND } \text{carry}) / (or' \text{ AND } \text{carry})$

$B \leftarrow o'; y \leftarrow y/8; o' < -o$

$\text{lsb}(a) = o'$

end for

if  $(-lo = \text{TRUE})$  then  $A \leftarrow \text{ADD}.(A, b)$

However, Menezes discloses instruction for performing encryption utilizing integer division which could be implemented as a "Computer For Loop Condition Statement". Menezes' iterative calculation of the encryption process is recited on [pg. 598, sect. 14.20]. Therefore given applicants "For Loop" instruction to be iterated through (e.g., carryout by the computer) the encryption process, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Meneze's "For Loop Condition Statement" by employing the well known feature of adding an additional iterative step (e.g.,  $(n+1)$ ) for which will enhance data encryption within a chip card [pg. 598, sect. 14.20].

16. As to claim 11, Menezes teaches a method where at the end, the following operations are performed : if  $(713 = \text{TRUE})$  then b -  $\text{CPL2n}(b)$  if  $(\sim y = \text{TRUE})$  then c -  $\text{CPL2}(c)$ . (i.e., ... teaches condition if...then... statement logic [pg. 598, sect. 14.20, line 3.1])

17. As to claims 12 and 13, although the teaching of Menezes discloses substantial features of the claim invention it does not disclose:

An electronic component comprising calculation means programmed to implement a method said calculation means comprising a central unit associated with a memory comprising several registers for storing the data a and b (claim 12).

A chip card comprising an electronic component according to Claim 12. (claim 13).

However, these features are well known in the art and would have been an obvious modification of the system disclosed by Menezes as introduced by Drexler. Drexler discloses:

An electronic component comprising calculation means programmed to implement a method said calculation means comprising a central unit associated with a memory comprising several registers for storing the data a and b (claim 12) (to provide encryption processing means using integer division on a chip card [abstract]).

A chip card comprising an electronic component according to claim 12, (claim 13). (to provide encryption processing means using integer division on a chip card [abstract]).

Therefore, given the teachings of Drexler, a person having ordinary skill in the art at the time of the invention would have recognized the desirability and advantage of modifying Menezes by employing the well known feature of chip card data encryption disclosed above by Drexler, for enhancing chip card security [abstract].

**Prior Art Made of Record**

The prior art made of record and not relied upon is considered pertinent to applicant's disclosure.

- a. Joye et al. (US Patent Publication No. 20040184604)

### **Contact Information**

Any inquiry concerning this communication or earlier communications from the examiner should be directed to BRYAN WRIGHT whose telephone number is (571)270-3826. The examiner can normally be reached on 8:30 am - 5:30 pm Monday -Friday.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, William Korzuch can be reached on (571) 272-7589. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/BRYAN WRIGHT/

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